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# Application Delivery Controller Implementation to Cyberlab Data Center

Bachelor's Thesis Information Technology

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Tiivistelmä		
Palveluiden häiriönsietokyky, saatavu vän tietoverkkoympäristöä. Siksi Appl yrityksille korvaamaton työkalu parant sä. Application Delivery Controller lait ominaisuuksia räätälöidä niiden tarjoa	ication Delivery Network aa juuri näitä ominaisuu teet ja ohjelmistot tarjoa	ing teknologiasta on tullut ksia palvelinympäristös- vat yrityksille monia eri
Opinnäytetyön tarkoituksena oli asent seen ja testata niiden kuormanjako-or työ antaa tietoa kuormanjaosta ja sen logioista sekä F5:stä.	ninaisuutta. Tehdyn työr	n lisäksi tämä opinnäyte-
Työssä käytettiin F5:n virtuaaliversiota Virtuaaliversiolla perehdyttiin F5-ohjel versiota ajettiin VMware Workstationis maan ulkoverkkoa sekä sisäverkkoa. seen, joka toimi palvelimena ja ulkove oli yhteensopivuusongelmia palvelime siitä, miten konfiguroida kuormanjakoa tiin palvelinhuoneeseen ja konfiguroiti den palvelimiin, oli VMware vSpheree figuroitiin käyttämään kuormanjakoa.	mistoon ennen laitteider sta ja koneen verkkokort Sisäverkon puoli oli kytk erkon puoli luokan verkko en ohjelmisto kanssa, mu a sekä SSL:ää. Laitteide in tukemaan etäyhteyttä	n saapumista. Virtuaali- teja käytettiin simuloi- tetty yhteen tietokonee- oon. Tässä järjestelyssä utta antoi kuitenkin kuvan en saavuttua ne asennet- Jotta F5:t saivat yhtey-
Lopuksi koko verkko testattiin tekemä	llä palvelunestohyökkäy	ksiä F5:een
Työ saatiin tehtyä pisteeseen, joka va saatiin toimimaan onnistuneesti.	adittiin kyseisenä ajanko	ohtana. Kuormanjako
Asiasanat		
ADC, kuormanjako, virtuaali server		

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Abstract		
Service resiliency, availability and security ments; that is why Application Delivery Ne ble tool for companies to better these qual cation Delivery Controller equipment and s options for companies to tailor them exactl	tworking technology has ities in their server envir software offer many diffe	s become an invalua- conments. The Appli-
The purpose of the bachelor's thesis was t server room and test the load balancing fe done, the thesis offers information on load ADCs and ADC specific technologies and	atures of the equipment balancers and load bala	. Along with the work
The work was done by using F5 BIG-IP vir ment. The virtual version was used to get f equipment arrived. The virtual version setu VMware workstation and using the host PC inside VLANs. The inside was connected to outside was connected to the classes netw issues with the server software but offered ancing and SSL. After the equipment arrive configured to support remote session. Befer ers, there was a need to make some chan ment. After the changes in the vSphere, the Last, the whole network was stress tested The work was completed to the point that with working load balancing. SSL and furth	amiliar with the F5 softwap consisted of the softw C's network cards to sime o a PC that was used as vork. This setup had sor some insight as how to ed they were installed in ore the F5 could get con ges in the VMware vSph e F5s were configured f by doing DDoS attacks was needed in the envir	vare before the vare running in nulate outside and s a server and the ne interoperability configure load bal- n the server room and nection to the serv- nere server environ- for load balancing. against the BIG-IP.
the future.		be implemented in
Keywords		
ADC, load balancing, virtual server		

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# LYHENNELUETTELO

ADC	Application Delivery Controller
ADN	Application delivery network
CPU	Central processing unit
DDoS	Distributed denial of service
DSwitch	Distributed Switch. Virtual switch used inside the VMware vSphere web application
Gb	Gigabit
НТТР	Hypertext transfer protocol
IP	Internet protocol
KVM	Kernel-based Virtual Machine. Full virtualization solu- tion for Linux on x86 hardware
MAC	Media access control
NAT	Network address translation
NIC	Network interface card
OS	Operating System
OSI-model	Open System Interconnection model
PC	Personal computer
SSL	Secure Socket Layer. Cryptographic protocol family
SFP	Small form-factor pluggable
SFP TCP	Small form-factor pluggable Transmission control protocol

## 1 INTRODUCTION

The idea for this bachelor's thesis came from the need for having strong server side protection for the Cyberlab game and a need to scatter the traffic to multiple servers. The machines that are specialized in this kind of behavior are the F5 BIG-IP equipment and software family. Two of the 4000 series machines were bought to do the work. They needed to be installed and configured for the lab environment.

Before the equipment arrived, the preliminary familiarization was done with the virtual version of the 11.6 software that was running in VMware Workstation. Even though the virtual version had full graphical support, it was not easy trying to configure it because of no earlier experience with F5 products. Luckily, the instructor had enough knowledge to get started with the configuration.

After getting more familiar with the software, a goal was set to make a virtual pool, and put a game server behind that so that it was possible to ping it from some client machine. After the ping test was completed, SSL-offloading was to be tested.

After the physical machines arrived, the real work started. The equipment was installed in the Cyberlab server room and configured relatively the same way as the virtual version. Because the virtual servers were running in VMware vSphere, there was a need to make some configuration changes to there.

After the vSphere was configured, the client got a connection to the servers. The project remained in this state for a while, until the whole Cyberlab network was stress tested.

## 2 CYBERLAB PROJECT

#### 2.1 Project overview

The idea of the project is to create an environment where game design, cyber security and data center knowledge work together to create a bigger environment, what is called Cyberlab. It tries to mimic what is done in real business server environments and at the same time, offers a learning platform for cyber security and tries to constantly evolve through different kinds of tests and analytics.

The learning part of the project is done with gamification. It consists of game platforms created in the data center, where the players have a hacking OS and a target network or system. The goal is to try and get information out of the target OS and score points as set by the game designers. This could mean that getting inside the target OS or finding some relevant information through hacking and so on, would score you points.

Because the platform is built this way, there is a possibility that someone might try to break into the system itself. This is can be beneficial if the game environment is built rigid enough not to get hacked easily. Thus it can be used to determine and analyze what kind of hacking attacks are being done. The environment is mainly used for the IT-branch of studies as of now, but can be made to accommodate the needs of other branches of study as well.

#### 2.2 Game environment architecture

The operating systems run in KVM host, which in turn is created in vSphere server cluster, as seen in figure 1. The users make connection to the servers through F5 BIG-IP ADC, which handles SSL-offloading, load balancing, access policies and general information security, illustrated in figure 2.

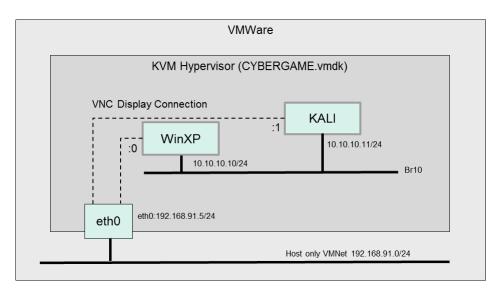


Figure 1. Game nest

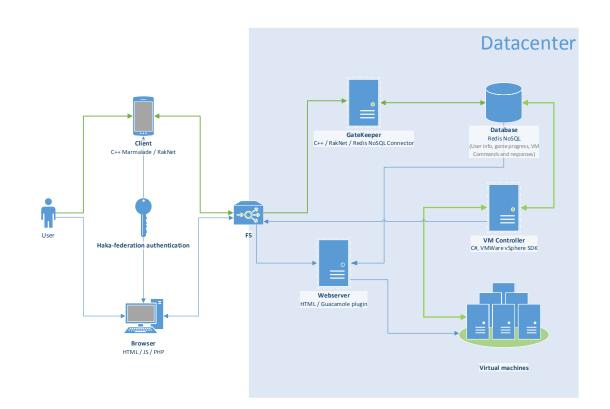


Figure 2. Cybergame connection architecture

## 3 LOAD BALANCING

Load balancing in network environment is a technique to distribute incoming traffic to multiple servers. It has become an important factor in today's internet architecture, as it is essential in server environment, where large numbers of users are trying to connect with the same content running in the servers. To provide all users continued access, there needs to be some sort mechanism to know when a server cannot accept any more connections and has to route the traffic to another server. The traffic can be distributed using many different styles.

#### 3.1 Evolution of load balancers

The first load balancing solutions for content servers used the cluster-IP, as seen in figure 3. This meant that even though the individual servers in a cluster had their own IP-address, the client would connect to the cluster-IP. The selection as to which server the client would connect, would be determined by which of the servers responded to the request the fastest. The OS in the servers was responsible for the load balancing. (Salchow 2012.)

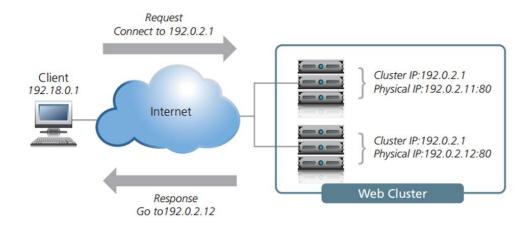


Figure 3. OS-based load balancing and cluster IP (Salchow 2012)

This kind of solution handled the load balancing well, if there was a small amount of clustered servers. When the server count rose to about 5 to 10, the communication between the servers started to impact the client connection. Also, if there happened to be any kind of disturbance in the solution used, it would impact the whole application and network that was running under it. (Salchow 2012.)

A step up from the OS based load balancing were the network-based load balancing hardware. These resided outside of the servers themselves and did not depend on a purpose built application, as seen in figure 4. (Salchow 2012.)

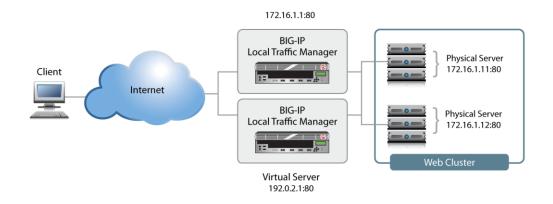


Figure 4. Network-based load balancing (Salchow 2012)

For the client to server connection the machines used virtual server addressing, similar as to how it is done by today's ADCs. The virtual server is created in the load balancers and is given its own IP-address to which clients are connecting. The load balancer then forwards the traffic to the servers, which are allocated to its use. Because the load balancers were now outside the servers themselves, it lessened the burden on the servers and the network. (Salchow 2012.)

High availability was strengthened simply because of the simpler solution. Also predictability became a new variant, because now that the equipment was outside the software itself, the traffic behavior could be monitored using response times, connection load and other variables. (Salchow 2012.) By analyzing this information, the load balancing could be customized for the needs of the company, for example, a more heavily utilized server does not receive as many connections as another server that has lesser load, based on server response time. The ADCs are heavily based on these designs, so they can be called the founding fathers of the ADCs. (Salchow 2012.)

## 3.2 Load balancing methods

## **Round Robin**

Round Robin load balancing sends the client request to all servers in a list like manner; first on the list gets it first, the second after that, and so on. The first server to announce it is available, becomes the active server for the client and is moved to the bottom of the list. (MacVittie 2009.)

This is easy to implement but has drawbacks. It assumes that all the servers on the list are up and have exactly the same load, storage and computing capacity. Round Robin has some variants that make it more viable solution to use. (MacVittie 2009.)

In Weighted Round Robin, the servers are assigned a weight number. The weight number indicates how the client requests are distributed among the servers according to assigned parameters, for example, 5:3:1 relation with three servers. This means that the server assigned with 5 gets the most of the requests, and so on. (MacVittie 2009.)

In Dynamic Round Robin, the system dynamically assigns the weight, based on a variety of parameters gathered from the real time data of the servers. (MacVittie 2009.)

## Random

As the name implies, Random load balancing method directs the traffic to a random server on the list using some kind of randomized algorithm and does not use any performance or status information to choose which server the traffic is forwarded to. Not very viable in today's high demand, high availability server environments, because of the simplicity of the design. (MacVittie 2009.)

## Fastest

Fastest option uses the server respond times to decide, as to which server to forward the traffic to; the fastest answering server is used for the client. It can be useful in situations, where the servers are scattered in different locations or networks. Its drawback is that a single response does not mean that it is the fastest after the response given and can so lead to congestion. (MacVittie 2009.)

#### Least Connections

Least Connections distributes the traffic to a server that has the least amount of connections. Works well, when the servers are all the same type; memory, computing capacity, and so on. It can cause congestion, because it does not know how long the connections are. For example, a single web page request versus multiple database queries. (MacVittie 2009.)

## Observed

Observed distributes the traffic to servers based on the combined values of least connections and fastest. The server with least connections and fastest response time is ranked the highest and gets the traffic. Can cause servers to be overloaded, because it does not know what kind of workload the server is going to be having after accepting the connections. (MacVittie 2009.)

#### Predictive

Uses the same ranking method as observed, but also checks, if the server performance is changing. If the server performance is improving, it gets the traffic and vice versa. (MacVittie 2009.)

## 4 APPLICATION DELIVERY CONTROLLER

ADCs are the next evolution of load balancers and SSL accelerators. They are platforms that do not only deliver load balancing features, but application acceleration, security and resiliency. They are very beneficial for businesses, which stretch over wide areas and are not confined in single location. In this kind of scenario, there could be multiple networks, protocols and restrictions working in each branch of the enterprise. The ADC can be made to accommodate the needs of every branch and optimize the traffic for each of them, while retaining the confidentiality, integrity and maximum availability of the data. Also if the enterprise has multiple data centers and one of them is experiencing congestion or goes down, the traffic can be forwarded to another data center even if it is located in another geological location.

In these kind of equipment there are a multitude of analytics and monitoring tools for optimizing the traffic and isolating different kinds of problems that might arise. For example, a service might be experiencing some lag for unknown reason. By using the tools to benchmark the service and system before, the problem can be found much faster by comparing the data with the benchmarked data and thus keep the system available and healthy for the users.

## 4.1 ADC used technologies

## Load balancing

Load balancing is widely used and essential for ensuring availability of services to customers and workers. See Chapter 3 for more information on load balancing techniques.

Inherent connection proxy and threat mitigation

Because the equipment resides in front of the main system, they add another layer of security for it. They can also, depending on the system, check the connection for any malcontent at any level of the OSI-model, as shown in figure 5.

		Attack	F5 Mitigation Technology
Increasing difficulty of attack detection	Application	OWASP Top 10 (SQL injection, XSS, CSRF, etc.), Slowloris, Slow POST, HashDos, GET floods	BIG-IP ASM: Positive and negative policy reinforcement, iRules, full proxy for HTTP, server performance anomaly detection
	Session	DNS UDP floods, DNS query floods, DNS NXDOMAIN floods, SSL floods, SSL renegotiation	BIG-IP LTM and BIG-IP GTM: High scale performance, DNS Express, SSL termination, iRules, SSL renegotiation validation
	Network	SYN floods, connection floods, UDP floods, PUSH and ACK floods, teardrop, ICMP floods, ping floods, and smurf attacks	BIG-IP AFM: SYN Check, default- deny posture, high-capacity connection table, full proxy traffic visibility, rate limiting, strict TCP forwarding

Figure 5. Overview how F5 ADCs cut the connection and can mitigate threats (Holmes 2013)

# Server task offloading

The ADCs can be used as a front end for the servers to offload different kinds of server intensive tasks.

SSL-offloading/acceleration is a technique to terminate encrypted data transmission before passing the traffic to the content server, as seen in figure 6. SSL-encryption and decryption are very intense for the CPU of the server, so making the ADC do the work for it helps ease the work load of the server (KEMP).

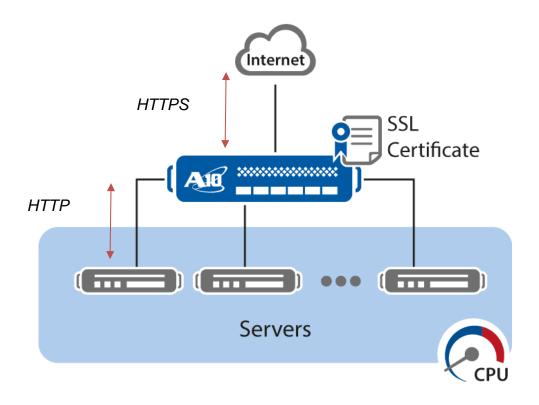


Figure 6. SSL-offloading (A10 Networks)

Before ADCs were integrated with this function, a dedicated CPU was handling the offloading task. This option was rather expensive, since there had to be one card for each server and its only function was to perform SSL. (KEMP.)

Then came the SSL accelerators, which operate in the same area as ADCs. They would be in front of the servers decrypting and encrypting the traffic. These proved to be very successful and were later paired with load balancers. (KEMP.)

TCP multiplexing is a technique to reuse non active TCP connections, as seen in figure 7. This helps to improve server capacity and performance since they do not have to close the ended connections and open new ones, but can use the open connections for other tasks or users. This can also be beneficial in design perspective as the servers can be made to take less connections than without multiplexing, thus having more resources for more important tasks or having less servers to run the hosted applications. (MacVittie 2008.)

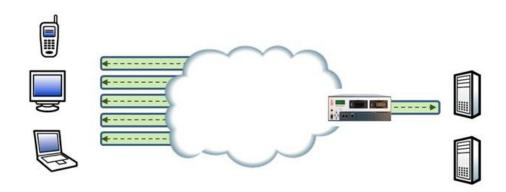


Figure 7. TCP multiplexing (MacVittie L 2008)

Advanced services

## Advanced firewall

Ensures data availability, integrity and confidentiality, using methods such as network session tracking, application awareness, attack mitigation and DDoS protection. (F5 Networks Inc. n.d.b.)

## Access policy manager

Differentiates users as to which resources they can access using user and group identities; such as location, device type and security posture. (F5 Networks Inc. n.d.a.)

## Application acceleration

Optimizes services and content by using HTTP-offloading and data caching, image downsizing, file size reduction and more. (F5 Networks Inc. n.d.c.)

## Secure web gateway

Works together with access policy manager to ensure maximum access to services and data, while keeping the assets healthy from malware and other web-borne threats using user access management, data integrity checks with Websense and configure actions taken after an attack or infection is noticed. (F5 Networks Inc. n.d.f.)

These are just a few advanced services that the companies' products offer. There are as many solutions as there are companies.

## 5 F5 NETWORKS

F5 is the market leader in ADN services (F5 Networks Inc. n.d.e). Its knowledge, innovation and resources differentiate it from its competitors (Fabbi&Lerner 2014).

F5 offers hardware and software solutions for companies' data, video and voice requirements. Its products include BIG-IP ADC hardware family, which offers load balancing, advanced firewalls, carrier grade NAT and many more solutions for companies to fully customize the machines for their needs. (F5 Networks Inc. n.d.e.). Their Silverline products offer cloud based solutions, like DDoS protection on demand and web application firewall (F5 Networks Inc. n.d.d.).

The company was founded in 1996 and has grown since to become a multimillion company boasting over 65 million dollars in shares and 2.3 billion dollars in assets. (F5.)

## 6 GETTING FAMILIAR WITH F5 SOFTWARE

Familiarization for the coming implementation of the equipment started with virtual version 11.6 of the F5 BIG-IP platform which ran over VMware Work-station. The first task was to get a server and the F5 see each other.

Before starting to configure the F5, there was a need to make changes and create new networks to VMware Workstation. The outside interface was bridged to the PC's main NIC (VMnet0), inside to a newly installed NIC (VMnet1) and the management interface used NAT to communicate with the host PC (VMnet2), as seen in figure 8.

Name	Туре	External Connection	Host Connection	DHCP	Subnet Address
VMnet0	Bridged	Intel(R) Ethernet Connectio	-	-	-
VMnet1	Bridged	Realtek PCIe GBE Family Co	-	-	-
VMnet8	Host-only	-	Connected	Enabled	192.168.126.0
VMnet2	NAT	NAT	Connected	-	10.0.0.0
VMnet3	Host-only	-	Connected	Enabled	192.168.247.0

Figure 8. WMnet configuration

After that, the created VMnets were paired with the interfaces of the virtual machine. Their MAC addresses needed to be checked so that they were the same in VMware and in the F5, as shown in figures 9 through 11.

Network Adapter Advanced Settings	
Incoming Transfer	
Bandwidth:	Unlimited 🔹
Kbps:	A V
Packet Loss (%):	0.0
Outgoing Transfer	
Bandwidth:	Unlimited 💌
Kbps:	A V
Packet Loss (%):	0.0
MAC Address	
00:0C:29:46:74:	2D Generate
ОК	Cancel Help

Figure 9. Checking MAC-address from VMware

[root@loc	alhost:INOPERATIVE:Standalone] config # ifconfig
eth0	Link encap:Ethernet HWaddr 00:0C:29:46:74:2D
	inet addr:10.0.0.50 Bcast:10.0.0.255 Mask:255.255.255.0
	inet6 addr: fe80::20c:29ff:fe46:742d/64 Scope:Link
	UP BROADCAST RUNNING MULTICAST MTU:1500 Metric:1
	RX packets:131 errors:0 dropped:0 overruns:0 frame:0
	TX packets:92 errors:0 dropped:0 overruns:0 carrier:0
	collisions:0 txqueuelen:1000
	RX bytes:15077 (14.7 KiB) TX bytes:7032 (6.8 KiB)

Figure 10. Making sure the F5 MAC-address is the same as in VMware

Virt	/irtual Machine Settings		
F	lardware Options		
	Device	Summary	
	Memory	8 GB	
	Processors	2	
	Hard Disk (SCSI)	104 GB	
	Hard Disk 2 (SCSI)		
	Network Adapter	Custom (VMnet2)	
	Network Adapter 2		
	t Network Adapter 3		
		Bridged (Automatic)	
	Display	Auto detect	

Figure 11. Network adapter alignment

Next step was to try and configure the software so, that it could get connection to a server. This proved to be rather difficult because of no earlier experience with F5 products, so a lot of trial and error was used, but finally the connection was made with the help from the instructor.

First, the inside VLAN, Self IPs and an outside route were created, as shown in figures 12 through 14.

Network » VLANs : VLAN List	» New VLAN			
General Properties				
Name	internal			
Description				
Тад				
Resources				
Interfaces	Untagged	1.4	Tagged	•

Figure 12. Creating VLAN

inside_1	192.168.0.1	255.255.255.0	internal
inside_2	192.168.0.2	255.255.255.0	internal
outside_1	10.0.0.2	255.255.255.0	external

Figure 13. Self-IP addresses

Network » Routes » New Route				
Properties				
Name	GW_outside			
Description				
Destination	0.0.0.0			
Netmask	0.0.0.0			
Resource	Use Gateway			
Gateway Address	IP Address			
MTU	0			

Figure 14. Configuring outside route

Then, a pool and virtual server were created for the server, as shown in figures 15 and 16.

Name	test_pool
Description	
Health Monitors	Active Available       Active     Available       /Common     (Common       http     (Common       HTTP_8080       http_head_f5       https       https       https_https_443
Resources	
Load Balancing Method	Round Robin 🔻
Priority Group Activation	Disabled •
New Members	New Node Node List Node Name: (Optional) Address: 172.16.20.10 Service Port: 80 HTTP  Add R:1 P:0 C:0 172.16.20.10 172.16.20.10 :80  Edit Delete

Figure 15. Creating test pool and adding a node

Name	test_vServer
Description	
Туре	Standard
Source	
Destination	Type:  Host Network Address: 10.69.48.50
Service Port	80 HTTP •
Notify Status to Virtual Address	
State	Enabled •

Figure 16. Creating virtual server

Last, the configuration was tested by pinging the inside VLAN server from an outside VLAN endpoint.

# 7 IMPLEMENTING BIG-IP TO THE CYBERLAB

The goal was to build an environment to test the load balancing of the BIG-IP. It consisted of an end machine, which simulated the user connecting from internet, the F5s, which load balanced the traffic, and the servers, which the user was trying to connect to, as seen in figure 17.

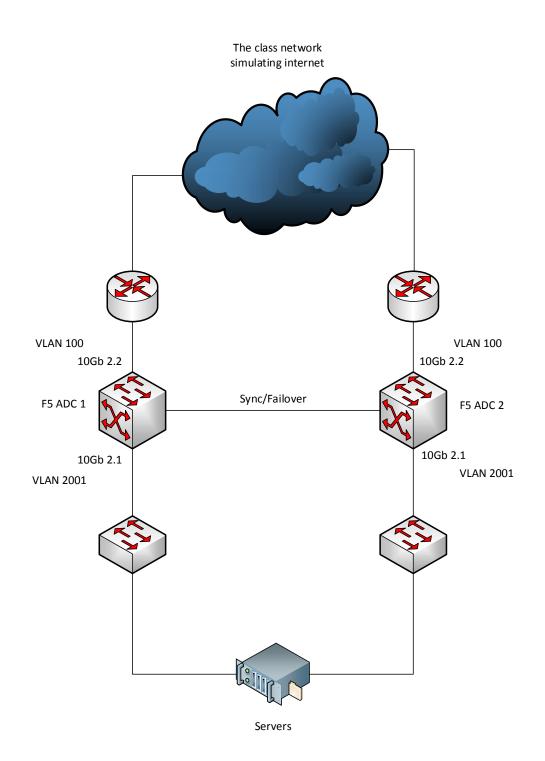


Figure 17. Project topology

## 7.1 Installation and initial configurations

The work started by installing the power supplies to the equipment and mounting them to the server rack. The installation was easy because the rack supports were a tool-less model. Next the default management IP-addresses were changed to temporary ones for access to the machines remotely.

The cables used were SFP copper and the interfaces that were used for the client and server connection were 10Gb ports. 2.2 ports were connected to Cisco ASR 900 routers which marked the outside and 2.1 ports were connected to a Dell switch, which forwarded the traffic to the inside network. Also the sync cable that came with the package was connected between the two F5 devices.

After getting control remotely, the machines were run through the initial configuration wizard which was prompted automatically when logging in the first time. There, the Self IPs, VLANs and outside routes were configured according to the laboratories IP-addressing.

VLANs were tagged; inside was tagged with 2001 and outside with 100.

## 7.2 Configuration sync

After the initial configuration was ran through, the machines were made to be in sync. This was done from the Device Management. First, a peer needed to be added. This needed the IP-address of the peer device and the configured credentials, as seen in figure 18.

	Device Management » Device Trust : Peer List				
F	Remote Device Credentials				
	Device IP Address				
	Administrator Username				
	Administrator Password				
	Cancel Retrieve Device Information				

Figure 18. Adding a peer device

Next step was to configure the network failover addresses, as shown in figure 19.

Device Management » Devices » cyberlab-adc1.ictlab.kyamk.fi				
tor → Properties Device Connectivity →				
Failover Unicast Configuration		Add		
✓ Local Address		\$ VLAN		
<u>10.69.48.25</u>	1026	Management Address		
172.16.0.2	1026	internal		
Delete				

Figure 19. Adding the needed failover addresses

Because the devices said they were in sync, they were left as they were. In reality this was not the case, because the network failover options and group syncing were not configured properly. The full syncing was done at a later time, as shown in figure 20, which resulted in the devices becoming properly synced, with the other device being the primary one and in active mode and the other secondary and in standby mode.

Device Management » Device Groups » SyncFailover_group				
🚓 🚽 Properties Failo	ver			
General Properties				
Name	SyncFailover_group			
Group Type	Sync-Failover			
Description				
Configuration: Advanced 🔻				
Members	Includes /Common cyberlab-adc1.ictlab.kyamk.fi cyberlab-adc2.ictlab.kyamk.fi	Available		
Automatic Sync				
Full Sync				
Maximum Incremental Sync Size (KB)	1024			
Update Delete				

Figure 20. Device group configuration

Finally, the configured device was made to sync its configuration to the group and overwrite the configuration of the other device, as seen in figure 21. This is a step that needs caution, because if careless, it can be made to the sync configuration of a device that is not wanted

🕁 🚽 Overview	Overview				
	-				
evice Groups					
≑ Name		Number of the second	f Devices	Group Type	Sync Type
SyncFailover_group	0	2	Sync-Faild	ver	Manual
device_trust_group	0	2	Sync-Only		Auto
Devices	Details				Show Advanced View
					Show Advanced view
♦ HA Status ▲ Name	1			Configu	ration Time
autoriat	-adc1.ictlab.kyamk.fi	(Self)	٩	1/19/2016	12:20:34
cyberlat	-adc2.ictlab.kyamk.fi		0	1/19/2016	12:20:34
•	ado2.ioudb.ityaiiit.ii				
•	ado <u>z.iodab.ityani</u> ti.i				

Figure 21. Syncing the main devices configuration to peer

# 7.3 Creating test virtual server

Before the sync, a virtual test server was made to see, if prior knowledge from the virtual version held true and the connection from a client to the server could be established.

The configuration started with creating a pool for servers and adding a host under it, as seen in figure 22.

Local Traffic » Pools : Pool List » New Pool				
Configuration: Basic 🔹				
Name	test_pool			
Description				
Health Monitors	Active Available       Active     Available       /Common     http       <<     /Common       http     <<       <     >>       https     https_443       https_head_f5			
Resources				
Load Balancing Method	Round Robin			
Priority Group Activation	Disabled			
New Members	New Node Node List  Node Name: (Optional)  Address: 192.168.0.10  Service Port: 80 HTTP  Add  R:1 P:0 C:0 192.168.0.10 192.168.0.10 :80  Edit Delete			

Figure 22. Creating a pool and adding a node

After the pool was configured, the next step was to create a virtual server for the pool, as seen in figure 23.

#### Local Traffic » Virtual Servers : Virtual Server List » New Virtual Server...

General Properties	
Name	VS_test
Description	
Туре	Standard
Source	
Destination	Type:  Host Network Address: 10.0.3.10
Service Port	80 HTTP V
Notify Status to Virtual Address	
State	Enabled V

Figure 23. Creating a virtual server

The configuration options that were in the previous figure are what mattered the most. Other options, except source address translation were left as they were, even though it is best practice to use a TCP-LAN-optimized profile for the client side protocol profile and have the virtual server enabled only on the external side.

The auto map option was used as source address translation. This option changes the clients' own IP to BIG-IP's self -IP address and forwards the traffic with that address as a source to the servers.

Last, the previously configured pool was added as a resource for the virtual server, as show in figure 24.

Resources		
	Enabled	Available
iRules	Up Down	<pre> /Common CYBERGAME_SVC_POLICY SlowLoris_Block .sys_APM_ExchangeSupport_OA_BasicAuth _sys_APM_ExchangeSupport_OA_NtlmAuth </pre>
Policies	Enabled	Available       <
Default Pool +	HelloWorld	<b>T</b>
Default Persistence Profile	None 🔻	
Fallback Persistence Profile	None 🔻	
Cancel Repeat Finished		

Figure 24. Adding pool as a resource to virtual server

## 7.4 VMware vSphere configurations

Next step was to make changes in the server cluster with vSphere web client, so that the traffic going for the servers behind BIG-IPs was isolated from the main link. This needed a creation of new distributed switch inside the server cluster, as shown in figure 25.

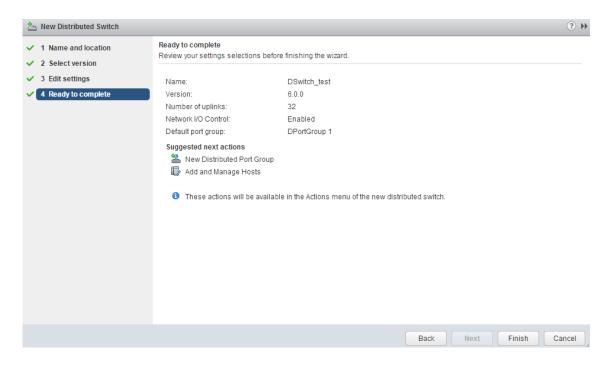


Figure 25. Creating a DSwitch in VMware web application

Next there was a need to configure new port group for the test lab, so that the DSwitch forwards the traffic to the right place. The inside was chosen to be in VLAN 2001, as seen in figure 26.

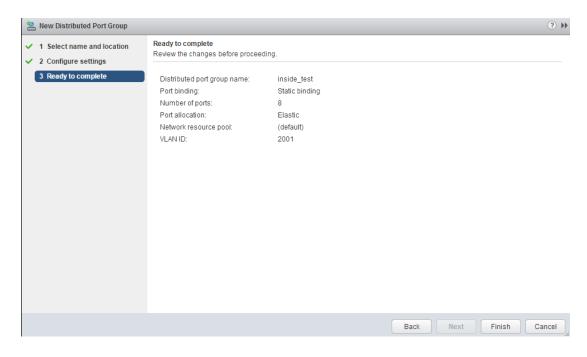


Figure 26. Port group configuration

After that, the needed server hosts were migrated under the newly created port group, as seen in figure 27.

HELLOWORLD1 - Edit	t Settings	(?) ₩
Virtual Hardware VM C	Options SDRS Rules vApp Options	
▶ 🔲 CPU	1 0	
▶ 🏧 Memory	1024 <b>v</b> MB <b>v</b>	
▶ 🚍 Hard disk 1	20 A GB V	
▶ 🛃 SCSI controller 0	LSI Logic Parallel	
Metwork adapter 1	Cyberlab INSIDE (CYBERLAB DSwit	
▶ iiiiiiiiiiiiiiiiiiiiiiiiiiiiiiiiiiii	Host Device Connect	
🕨 🖶 USB controller	USB 2.0	
▶ 🛄 Video card	Specify custom settings	
VMCI device		
▶ 🝥 SATA controller 0		
<ul> <li>Other Devices</li> </ul>		

Figure 27. Adding host to port group

After the configurations were done, ping command was used to test if the server node was up and active, as shown in figure 28.

[root@cy	yberlab-adc1:Active:In Sync] config # ping 172.16.30.10
PING 172	2.16.30.10 (172.16.30.10) 56(84) bytes of data.
64 bytes	s from 172.16.30.10: icmp_seq=1 ttl=64 time=0.578 ms
64 bytes	s from 172.16.30.10: icmp_seq=2 ttl=64 time=0.807 ms
64 bytes	s from 172.16.30.10: icmp_seq=3 ttl=64 time=0.912 ms

Figure 28. Ping test from F5 to server

#### 7.5 Load balance test

After the environment was tested and was deemed ready for further experimenting, a goal was set to create 3 different web servers and see how the F5 would load balance the traffic from the clients to those.

The web servers were made with Linux and a simple html index page was used to demonstrate the case. Three of these pages were created with different index pages and IP-addresses, as seen in figures 29 and 30.

<html></html>			
<pre><head></head></pre>			
<title>TEST&lt;/td&gt;&lt;td&gt;SERV&lt;/td&gt;&lt;td&gt;1</title>			
<body></body>			
<h1>Test_serv</h1>	1 <td>&gt;</td> <td></td>	>	
<pre></pre>			

Figure 29. Test server index page

iface eth0 inet static address 172.16.20.10
address 172.16.20.10
netmask 255.255.0.0
gateway 172.16.0.1

Figure 30. Test server IP-addresses

After they were created, their interfaces needed to be brought down and up for the new IP-addresses to take effect. Then, they were migrated under the DSwitch inside in vSphere, so that they would get connection to the Big-IP.

Next, a pool and virtual server was created for the environment, as shown in figures 31 through 34.

Local Traffic » Pools : Pool List » HelloWorld				
🔅 👻 Properties	s Members		Statistics 💌	
General Properties		-		
Name	HelloWorld			
Partition / Path	Common			
Description				
Availability	🔷 Offline (I	Enabled) The	children pool member(s	s) are down
Configuration: Basic	•			
Health Monitors	Act /Common http	ive	Available /Common gateway_icmp http_head_f5 https https_443	•
Update Delete				

Figure 31. Pool for the web servers

Local Traffic » Pools : Pool List » HelloWorld				
🚓 🚽 Properties	Memb	ers	Statistics	
Load Balancing				
Load Balancing Method		Round Rob	in	•
Priority Group Activation Disabled				
Update				
Current Members				
Status 🗢 Memi	ber			
172.16.	20.10:80			
172.16.	20.20:80			
172.16.20.30:80				
Enable Disable Remove				

Figure 32. Servers used in the pool

Local Traffic » Virtual Servers : Virtual Server List » HelloWorld_vSRV				
🕁 🚽 Properties	Resources Security - Statistics >			
General Properties				
Name	HelloWorld_vSRV			
Partition / Path	Common			
Description				
Туре	Standard			
Source	0.0.0/0			
Destination	Type:  Host Network Address: 193.167.58.248			
Service Port	80 HTTP V			
Notify Status to Virtual Add	ress 🗹			
Link	None			
Availability	Offline (Enabled) - The children pool member(s) are down			
Syncookie Status	Off			
State	Enabled V			

Figure 33. Virtual server configuration

Local Traffic » Virtual Servers : Virtual Server List » HelloWorld_vSRV						
🗱 🚽 Properties	Resou	ırces	Security		Statistics	
Load Balancing						
Default Pool		HelloWorld			•	
Default Persistence	Profile	None	T			
Fallback Persistence	e Profile	None	T			
Update						

Figure 34. Adding pool to the virtual server

Finally, the system was tested by trying to connect to the virtual server IP and see if the ADC balanced the traffic to each node as the traffic of the servers increased; and it worked.

## 8 NETWORK STRESS TEST

After the BIG-IPs and the network were configured, a stress test of the whole environment was done by launching different kinds of DDoS attacks against the BIG-IP, as shown in figure 35.

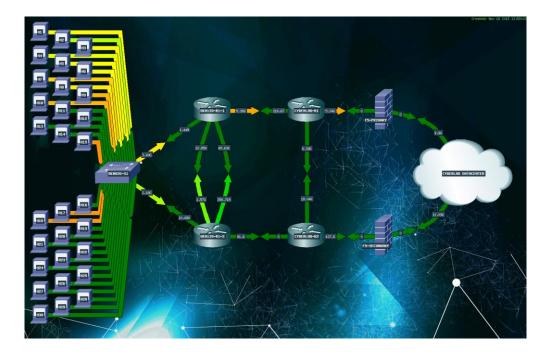


Figure 35. Stress test network

The main attacks used were TCP SYN flood and Slowloris. SYN flood uses TCP state retention to its advantage when a new connection is being established. By sending large amount of SYN packets against the target machine, it stops the target's ability to establish new legitimate connections. (Eddy 2007.)

Slowloris uses HTTP connection timer to its advantage by sending HTTP GET request very slowly. For example, a server has 200 second timer before it closes the connection if the client does not finish the GET request. By sending a partial GET and then feeding bogus information before the timer expires, the connection stays open. Sending large amount of these requests a hacker can overwhelm the targets ability to take any legitimate connections. (Muscat 2013.)

The testing spanned two days with one day done only with students. The second day a F5 consultant came to help with the configuration of the BIG-IP.

At start, BIG-IP held fairly well against the attacks without specific configurations, but as the severity of the attacks rose, the BIG-IP started showing extreme lag and poor availability of resources.

First module used against the attacks was the DoS protection profile and under it Application Security. Also Source IP-based Rate Limiting and URL-Based Rate Limiting, were enabled as shown in figure 36.

Security » DoS Protection : DoS	Profiles » DoS Profile Properties		
👷 👻 DoS Profile Properties			
eneral Configuration			
Profile Name	dos		
Partition / Path	Common		
Application Security	Enabled		
Protocol Security (DNS)	Enabled		
Protocol Security (SIP)	Enabled		
Application Security			
Trigger iRule	Enabled		
FPS-based Anomaly			
Operation Mode	Blocking V		
Prevention Policy	<ul> <li>Source IP-Based Client Side Integri</li> <li>URL-Based Client Side Integrity De</li> <li>Site-wide Client-Side Integrity Defer</li> <li>Source IP-Based Rate Limiting</li> <li>URL-Based Rate Limiting</li> <li>Site-wide Rate Limiting</li> </ul>	fense	
	Note: Blocked requests will be rejected	at the TCP Layer by this prevention policy	
IP Detection Criteria	TPS increased by	500 %	
Set default criteria	TPS reached Minimum TPS Threshold for detection	200 transactions per second 40 transactions per second	
URL Detection Criteria Set default criteria	TPS increased by TPS reached Minimum TPS Threshold for detection	500       %         1000       transactions per second         200       transactions per second	
Prevention Duration Set default duration		onds onds	

Figure 36. DoS Profile configuration

After the settings were done, the BIG-IP was able to forward traffic to the servers even though 4Gb of traffic went to it.

The next attack used was Slowloris. This stopped the BIG-IP service. Many configuration changes were tried against the attack but none were able to get the service running, until Single Endpoint Flood and Sweep values were changed, as shown in figure 37. It remained unclear this was the correct way of mitigating the attacks because while it stopped the attack it also stopped the monitoring tool used in the test. These were the last configuration changes done the first day.

Security » DoS Protection : I Properties	Device Configuration » Single Endpoint Flood
Attack Type	Single Endpoint Flood
Detection Threshold PPS	Specify • 4000
Default Internal Rate Limit	Specify • 8000
Packet Type	Selected     Available       Any IPv4     Any ICMP (IPV4)       Any UDP (IPv4)     Any UDP (IPv6)       TCP SYN without ACK (IPv4)     CP SYN without ACK (IPv6)

Figure 37. Single Endpoint Flood and Sweep configurations

The next day of testing started the same way as the first, with rising attack severity, but with the help of the F5 consultant, the BIG-IP was able to mitigate all of the attacks at the end. All in all the BIG-IP was deemed very effective against the attacks and fully serves its purpose in the Cyberlab environment.

## 9 CONCLUSIONS

The work was completed to the point that was needed at that time in the environment; with working load balancing. It provided a good learning experience as to how the ADC distributes the traffic and how to get traffic running through it.

The work done with the vSphere provided a small experience as how to configure virtual switches and distribution port groups and how to use the web application in general.

SSL-offloading was tested with the virtual version, but did not yield conclusive evidence that it worked properly. It needs to be tested with the hardware and implemented when the need arises.

The F5 iRule module was tested but because of no programming experience, was not used to its full potential. Used properly, it could be very beneficial for the Cyberlab project.

The primary modules that were used in the ADC were local traffic, network and device management, so there is much more to discover and use to optimize the F5 BIG-IP for the needs of the project and to learn of the machine in general. The F5 offers many more bachelor's thesis ideas such as IPv4 to IPv6, firewall, packet manipulation.

The work proved very challenging at start because of no previous experience with F5 products. The virtual testing environment was not very effective because of the interoperability issues with the server software and thus created more problems at the start. Also because the F5 has so many modules and functions, a lot of time was consumed trying to find the right commands for the right functions. A10 Networks Inc. n.d. SSL Offload. Available at:

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Figure 1. Game nest.

Figure 2. Cybergame connection architecture.

Figure 3. *OS-based load balancing and cluster IP.* Salchow, K Jr. 2012. Load Balancing 101: The Evolution of Application Delivery Controllers. F5 Network Inc. 4.23.2012. Available at: <u>https://f5.com/zh/resources/white-papers/load-balancing-101-the-evolution-to-application-de [Accessed: January 10 2016].</u>

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Figure 5. Overview how F5 ADCs cut the connection and can mitigate threats. Holmes, D. 2013. Mitigating DDoS Attacks with F5 Technology. F5 Networks Inc. 1.31.2013. Available at: <u>https://f5.com/resources/white-papers/mitigating-</u> <u>ddos-attacks-with-f5-technology</u> [Accessed: February 8 2016].

Figure 6. *SSL-offloading.* A10 Networks Inc. n.d. SSL Offload. Available at: <u>https://www.a10networks.com/products/ssl-offload</u> [Accessed: January 25 2016].

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Figure 8. VMnet configuration.

Figure 9. Checking MAC-address from VMware

Figure 10. Making sure the F5 MAC-address is the same as in VMware

Figure 11. Network adapter alignment.

Figure 12. Creating VLAN.

Figure 13. Self-IP addresses.

Figure 14. Configuring outside route.

- Figure 15. Creating test pool and adding a node.
- Figure 16. Creating virtual server.
- Figure 17. Project topology.
- Figure 18. Adding a peer device.
- Figure 19. Adding the needed failover addresses.
- Figure 20. Device group configuration.
- Figure 21. Syncing the main devices configuration to peer.
- Figure 22. Creating a pool and adding a node.
- Figure 23. Creating a virtual server
- Figure 24. Adding pool as a resource to virtual server.
- Figure 25. Creating a DSwitch in VMware web application.
- Figure 26. Port group configuration.
- Figure 27. Adding host to port group.
- Figure 28. Ping test from F5 to server.
- Figure 29. Test server index page.
- Figure 30. Test server IP-addresses.
- Figure 31. Pool for the web servers.
- Figure 32. Servers used in the pool.
- Figure 33. Virtual server configuration.
- Figure 34. Adding pool to the virtual server.
- Figure 35. Stress test network.
- Figure 36. DoS Profile configuration.
- Figure 37. Single Endpoint Flood and Sweep configurations.